Guaranteed Throughput and Best Effort Streams in a Single Network on Chip Model

GDR!

ZUT 2009

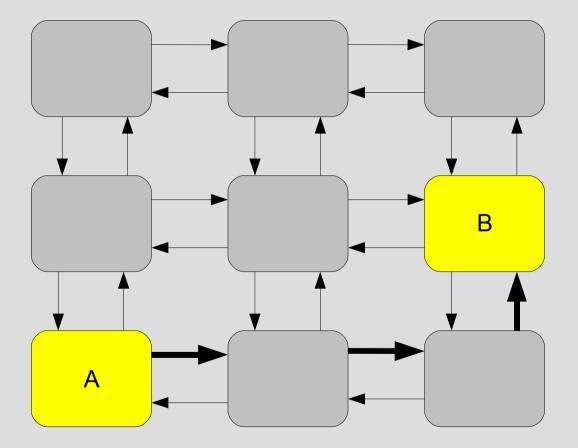
Problem description

Network on Chip Balance the load Save queue length Reduce hotspots Quality of Service

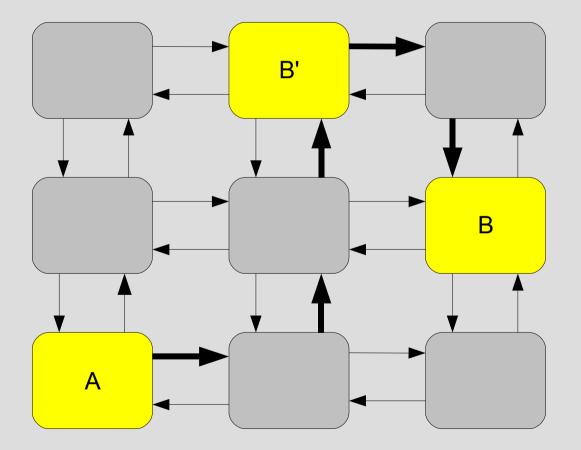
Previous work

Variants of XY algorithm Valiant algorithm Random XY/YX Described as near-optimal!? Classical algorithms aren't QoS-aware

XY



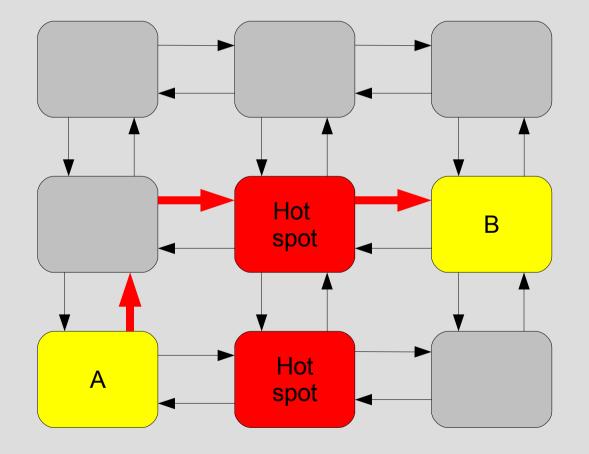
Valiant



My proposition (PARouting)

Different routes for different priorities Best Effort packets don't need minimal routes Local algorithm Could improve balancing Not aware of global load spread Could make balancing ever worse!

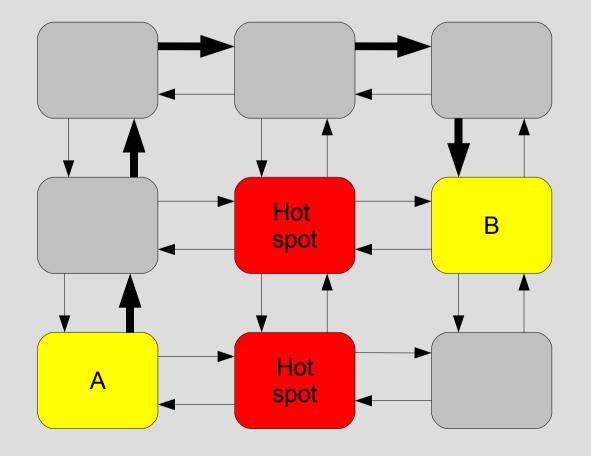
My proposition (PARouting)



Lo pri

Hi pri

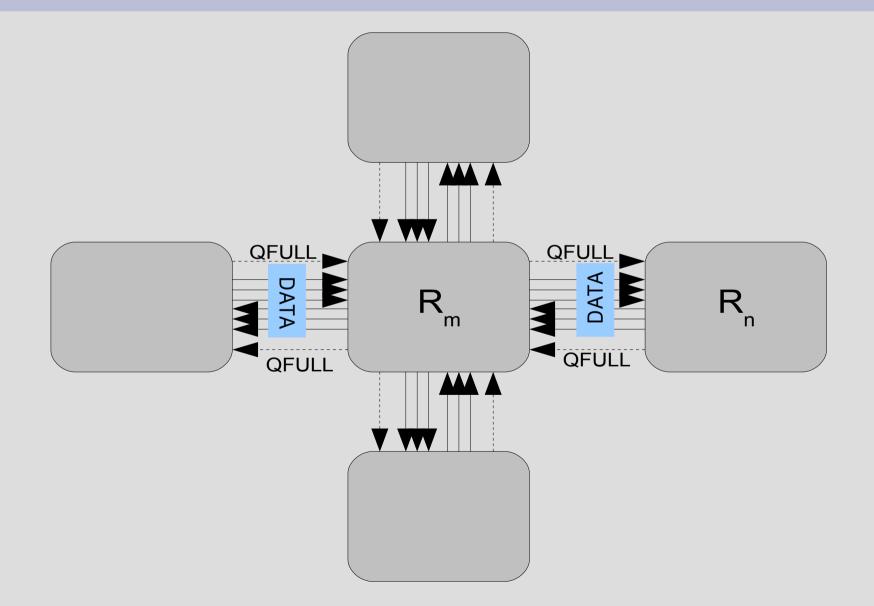
My proposition (PARouting)



Lo pri

Hi pri

PARouting – Q_{FULL} signal



PARouting - N_{bad}

Extra field in a packet – N_{bad} 2-5 bits

If N_{bad} non-zero: "Bad" hop may be performed If "Bad" hop performed: N_{bad} is decreased

PARouting - N_{bad}

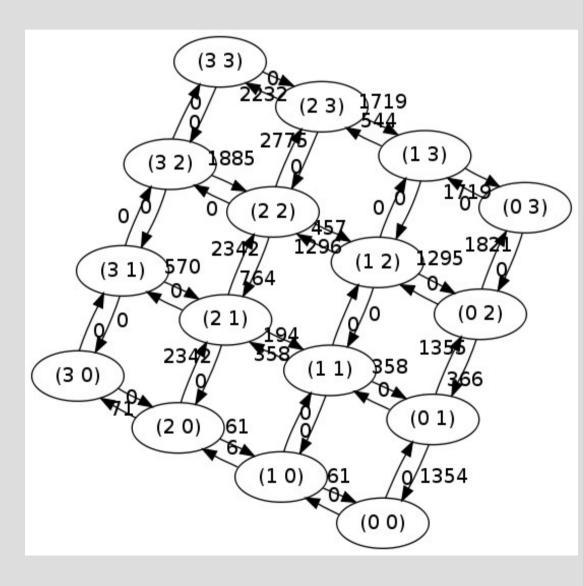
 N_{bad} initialized depending on priority Depending on mesh size Lower priority packets get higher N_{bad} Highest priority packets have $N_{bad} = 0$

PARouting – selecting direction

Calculate D_{XY} and D_{YX} directions for XY/YX If any of routers at D_{xy} , D_{yy} don't emit QFULL Send the packet there Else If $N_{bad} > 0$ Send packet to a non-minimal direction Else $(N_{bad} = 0)$ Send packet to D_{yy}

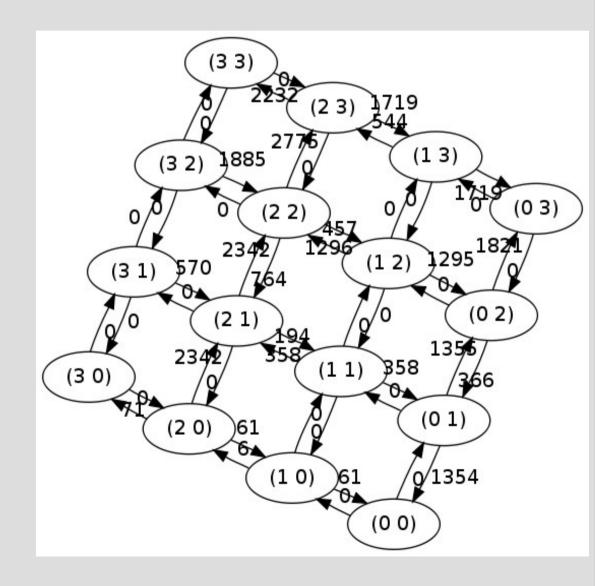
Simulation

Number of flits passed Queue lengths Hops Average trip time

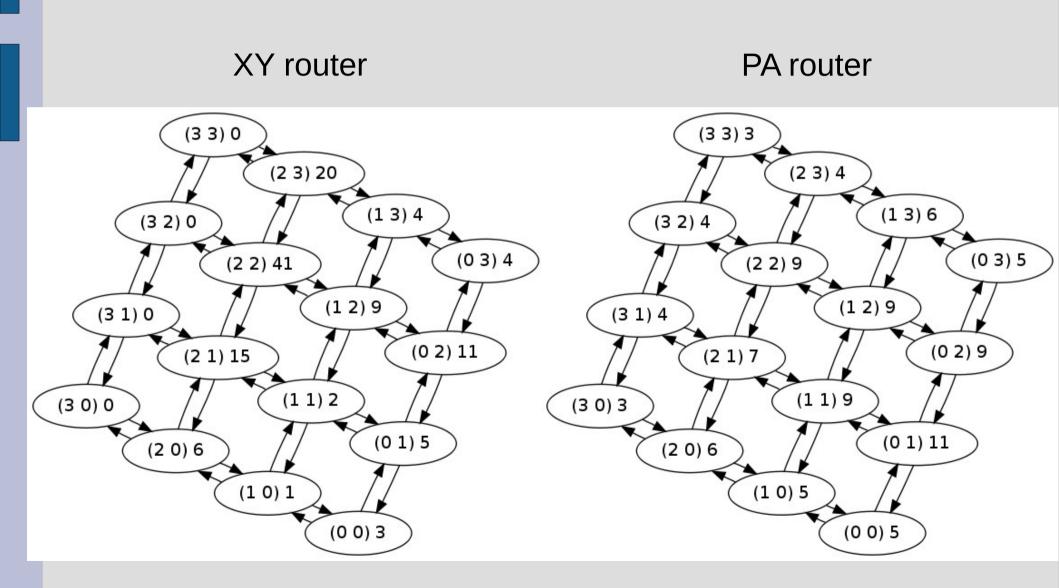


Simulation

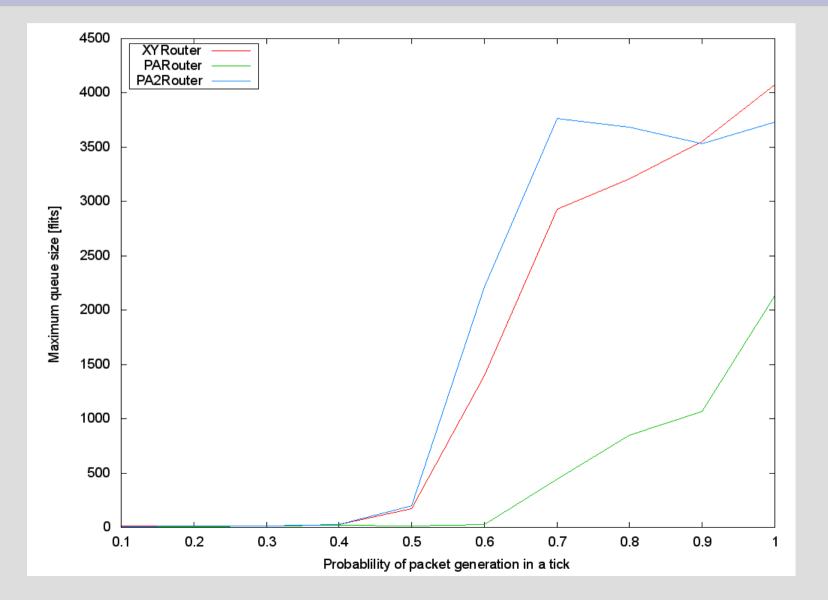
Python GraphViz Gnuplot



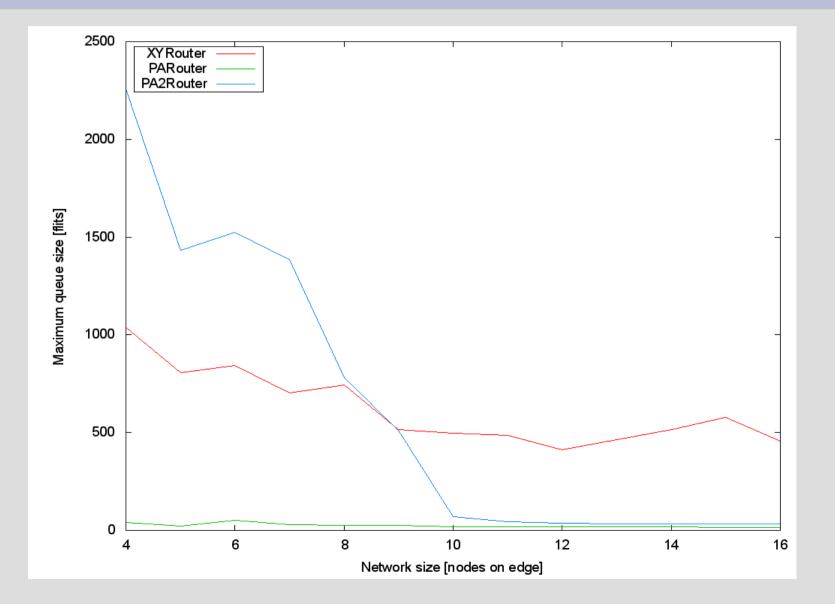
Results: reducing hotspots



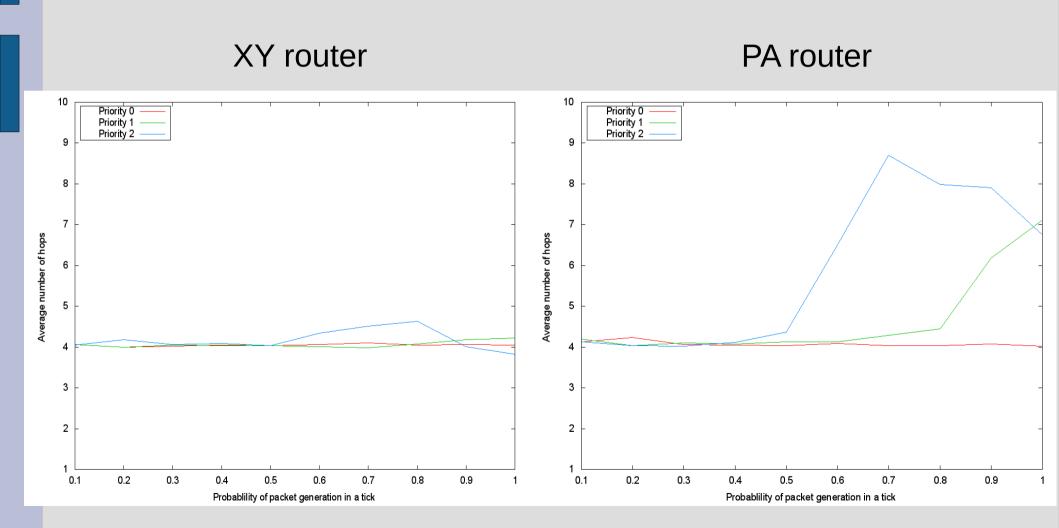
Results: queue sizes



Results: queue sizes



Results: hops



Summary

Better load distribution between nodes Smaller queue sizes At the cost of latency But not of high-priority packets